

last time All rings are commutative and have 1

## Categories

• objects ( $\mathcal{O}$ )

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• arrows  $A \xrightarrow{f} B$  with source  $A$  and target  $B$

such that:

• Arrows  $A \xrightarrow{f} B$  and  $B \xrightarrow{g} C$  can be composed to make an arrow  $A \xrightarrow{gf} C$  and composition is associative

• For every object  $A$ , there is a special arrow  $1_A \in \text{Hom}_{\mathcal{O}}(A, A)$

$$A \xrightarrow{1_A} A \xrightarrow{f} B = A \xrightarrow{f} B, \quad B \xrightarrow{g} A \xrightarrow{1_A} A = B \xrightarrow{g} A$$

Notation: the collection of all arrows  $A \rightarrow B$  is  $\text{Hom}_{\mathcal{O}}(A, B)$

## Examples

- 1) Set : objects are all sets, arrows are all functions
- 2) Grp : groups and group homomorphisms
- 3) Ab : abelian groups and homomorphisms
- 4) Ring : rings and ring homomorphisms
- 5) Top : topological spaces and continuous functions
- 6) R-mod : R-modules and R-module homomorphisms  
also known as:  $\text{Mod}(R)$  (vs  $\text{mod}(R) = \text{fg R-modules}$ )

## Another example

$X$  partially ordered set

We can regard  $X$  as a category

objects  $x \in X$

arrows if  $x \leq y \Rightarrow \text{Hom}_{\mathcal{C}}(x, y)$  singleton

if  $x \not\leq y \Rightarrow \text{Hom}_{\mathcal{C}}(x, y) = \emptyset$

## Special types of arrows:

•  $f \in \text{Hom}_{\mathcal{C}}(A, B)$  is an iso  $\equiv$  isomorphism if there exists  $g \in \text{Hom}_{\mathcal{C}}(B, A)$  such that  $gf = 1_A$ ,  $fg = 1_B$

•  $f \in \text{Hom}_{\mathcal{C}}(B, C)$  is monic or a monomorphism if

$$A \begin{array}{c} \xrightarrow{g_1} \\ \xrightarrow{g_2} \end{array} B \xrightarrow{f} C \quad fg_1 = fg_2 \Rightarrow g_1 = g_2$$

•  $f \in \text{Hom}_{\mathcal{C}}(A, B)$  is epi or an epimorphism if

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{g_1} \\ \xrightarrow{g_2} \end{array} C \quad g_1f = g_2f \Rightarrow g_1 = g_2$$

## Functors $\mathcal{B}, \mathcal{D}$ categories.

A covariant functor  $F: \mathcal{B} \rightarrow \mathcal{D}$  is a mapping that assigns

• to each object  $C$  in  $\mathcal{B}$ , an object  $F(C)$  in  $\mathcal{D}$

• to each arrow  $f \downarrow$  in  $\mathcal{B}$ , an arrow  $\downarrow F(f)$  in  $\mathcal{D}$

$$\begin{array}{ccc} & A & F(A) \\ & \downarrow f & \downarrow F(f) \\ & B & F(B) \end{array}$$

such that:

$$\textcircled{1} F(1_A) = 1_{F(A)}$$

$$\textcircled{2} F(fg) = F(f)F(g) \quad \text{for all composable arrows } fg$$

A Contravariant functor  $F: \mathcal{B} \rightarrow \mathcal{D}$  flips all arrows:

• to each object  $C$  in  $\mathcal{B}$ , an object  $F(C)$  in  $\mathcal{D}$

• to each arrow  $f \downarrow$  in  $\mathcal{B}$ , an arrow  $\uparrow F(f)$  in  $\mathcal{D}$

$$\begin{array}{ccc} & A & F(A) \\ & \downarrow f & \uparrow F(f) \\ & B & F(B) \end{array}$$

such that:

$$\textcircled{1} F(1_A) = 1_{F(A)}$$

$$\textcircled{2} F(fg) = F(g)F(f)$$

$$\begin{array}{ccc} & A & F(A) \\ & \downarrow g & \uparrow F(g) \\ & B & F(B) \\ & \downarrow f & \uparrow F(f) \\ & C & F(C) \end{array} \rightsquigarrow$$

A contravariant functor  $\mathcal{C} \rightarrow \mathcal{D}$

a covariant functor  $\mathcal{C}^{\text{op}} \rightarrow \mathcal{D}$

where  $\mathcal{C}^{\text{op}}$  is a category obtained from  $\mathcal{C}$  by flipping all the arrows

so an arrow  $f \downarrow$  in  $\mathcal{C}$   $\equiv$  an arrow  $\uparrow f$  in  $\mathcal{C}^{\text{op}}$

### Examples of Functors

#### ① Forgetful functor

Grp  $\rightarrow$  Set : forgets the group structure

R-mod  $\rightarrow$  Set : forgets the R-mod structure

#### ② Identity Functor $\mathcal{C} \rightarrow \mathcal{C}$

#### ③ Localization Fix ring $R$ , multiplicative subset $W \ni 1$

$W^{-1} : R\text{-mod} \rightarrow W^{-1}R\text{-mod}$

$M \longmapsto W^{-1}M$

$f \downarrow \longmapsto W^{-1}f \downarrow$   $\Rightarrow \frac{m}{w} \downarrow$   
 $N \qquad \qquad \qquad W^{-1}N \qquad \qquad \frac{f(m)}{w}$



Special types of functors:  $\mathcal{C}, \mathcal{D}$  locally small.  $F: \mathcal{C} \rightarrow \mathcal{D}$  is

faithful  $\text{Hom}_{\mathcal{C}}(A, B) \longrightarrow \text{Hom}_{\mathcal{D}}(F(A), F(B))$  all injective

full  $\text{Hom}_{\mathcal{C}}(A, B) \longrightarrow \text{Hom}_{\mathcal{D}}(F(A), F(B))$  all surjective

fully faithful full and faithful

embedding fully faithfully and injective on objects

A subcategory  $\mathcal{C}$  of  $\mathcal{D}$  is full if for all objects  $A, B$  in  $\mathcal{C}$   
 $\text{Hom}_{\mathcal{C}}(A, B) = \text{Hom}_{\mathcal{D}}(A, B)$ .

Equivalently, the inclusion functor  $\mathcal{C} \rightarrow \mathcal{D}$  is full

Examples

①  $\text{Ab}$  is a full subcategory of  $\text{Group}$

② The forgetful functor  $\mathbb{R}\text{-mod} \rightarrow \text{Set}$  is faithful, not full.

$F, G: \mathcal{C} \rightarrow \mathcal{D}$  functors

A natural transformation  $\eta: F \Rightarrow G$  is a mapping

object  $C \in \mathcal{C} \longmapsto \eta_C \in \text{Hom}_{\mathcal{D}}(F(C), G(C))$

such that

$$\text{for all arrows } f \downarrow \begin{matrix} A \\ B \end{matrix} \text{ in } \mathcal{C} \rightsquigarrow \begin{array}{ccc} F(A) & \xrightarrow{\eta_A} & G(A) \\ F(f) \downarrow & & \downarrow G(f) \\ F(B) & \xrightarrow{\eta_B} & G(B) \end{array} \text{ commutes}$$

$\eta$  is a natural isomorphism if  $\eta_A$  is an iso for all  $A$ .

$\text{Nat}(F, G) :=$  natural transformations from  $F$  to  $G$

can construct a functor category  $\mathcal{D}^{\mathcal{C}}$  with:

- objects all functors  $\mathcal{C} \rightarrow \mathcal{D}$
- arrows all natural transformations between such functors

Example  $\text{Id}: \text{Grp} \rightarrow \text{Grp}$ ,  $\text{ab}: \text{Grp} \rightarrow \text{Grp}$   
 $G \mapsto G^{\text{ab}} = G/[G, G]$

the mapping  $\pi: \text{Id} \Rightarrow \text{ab}$  is a natural transformation:

$$G \mapsto \begin{array}{c} G \\ \downarrow \pi_G \\ G^{\text{ab}} \end{array} \text{ quotient map}$$

$$\begin{array}{ccc} G & \xrightarrow{\pi_G} & G^{\text{ab}} \\ f \downarrow & & \downarrow f^{\text{ab}} \\ H & \xrightarrow{\pi_H} & H^{\text{ab}} \end{array}$$

for all  $G \xrightarrow{f} H$   
 group homomorphism.

Hom functors  $\mathcal{C}$  locally small category

the covariant functor

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(A, -) : \mathcal{C} & \longrightarrow & \text{Set} \\ B & \longmapsto & \text{Hom}_{\mathcal{C}}(A, B) \\ \begin{array}{c} B \\ f \downarrow \\ C \end{array} & \longmapsto & \begin{array}{ccc} \text{Hom}_{\mathcal{C}}(A, B) & \ni & g \\ \downarrow & & \downarrow \\ \text{Hom}_{\mathcal{C}}(A, C) & \ni & f \circ g \end{array} \end{array}$$

$$\text{Hom}_{\mathcal{C}}(A, f) =: f_*$$

$$f_*(g) = fg \quad \begin{array}{ccc} A & \xrightarrow{g} & B \\ & \searrow & \downarrow f \\ & & C \end{array}$$

the contravariant

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(-, B) : \mathcal{C} & \longrightarrow & \text{Set} \\ A & \longmapsto & \text{Hom}_{\mathcal{C}}(A, B) \\ \begin{array}{c} A \\ f \downarrow \\ C \end{array} & \longmapsto & \begin{array}{ccc} \text{Hom}_{\mathcal{C}}(A, B) & \ni & \\ \uparrow & & \uparrow \\ \text{Hom}_{\mathcal{C}}(C, B) & \ni & g \end{array} \end{array}$$

$$\text{Hom}_{\mathcal{C}}(f, B) =: f^*$$

$$f^*(g) = gf \quad \begin{array}{ccc} A & \xrightarrow{f} & C \\ \vdots \downarrow & \swarrow g & \\ B & & \end{array}$$

Yoneda Lemma  $\mathcal{C}$  locally small category

$F: \mathcal{C} \rightarrow \text{Set}$  covariant functor

Fix an object  $A$  in  $\mathcal{C}$

there is a bijection

$$\text{Nat}(\text{Hom}_{\mathcal{C}}(A, -), F) \xrightarrow[\cong]{\delta} F(A)$$

this correspondence is natural in both  $A$  and  $F$ .

Proof  $\varphi: \text{Hom}_{\mathcal{C}}(A, -) \Rightarrow F$

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(A, A) & \xrightarrow{\varphi_A} & F(A) \\ \downarrow \text{id}_A & \longmapsto & \varphi_A(\text{id}_A) = \mu \end{array}$$

Set  $\delta(\varphi) := \varphi_A(\text{id}_A) = \mu$ . Why is  $\delta$  injective/surjective?

Fix an arrow  $A \xrightarrow{f} X$ .  $\varphi$  natural transformation gives

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(A, A) & \xrightarrow{\text{Hom}_{\mathcal{C}}(A, f)} & \text{Hom}_{\mathcal{C}}(A, X) \\ \downarrow \varphi_A & \Downarrow & \downarrow \varphi_X \\ F(A) & \xrightarrow{F(f)} & F(X) \end{array}$$

$\begin{array}{ccc} \text{id}_A & \longmapsto & f \\ \downarrow & & \downarrow \\ \mu & \longmapsto & \boxed{\phantom{f}} \end{array}$

$$\begin{aligned} \boxed{\phantom{f}} &= \varphi_X \circ \text{Hom}_{\mathcal{C}}(A, f)(\text{id}_A) = \varphi_X(f) \quad \curvearrowright \\ &= F(f) \circ \varphi_A(\text{id}_A) = F(f)(\mu) \quad \hookrightarrow \end{aligned}$$

so:

- $\varphi$  is completely determined by  $\varphi_A(1_A)$ :

$$\varphi_x(f) = F(f)(\varphi_A(1_A))$$

so the values of  $\varphi$  are set for any  $x, f$  from  $\varphi_A(1_A)$

$\therefore$  different choices of  $u = \varphi_A(1_A) \Rightarrow$  different  $\varphi$

$\Rightarrow \delta$  is injective

$\therefore$  given any  $u \in F(A)$ , setting  $\varphi_x(f) = F(f)(u)$

gives a natural transformation  $\varphi$  with  $\delta(\varphi) = u$ .

$\Rightarrow \delta$  is surjective