Sidenote: Universal poperties

We've all seen constructions satisfying a universal property. these are best explained by examples:

Examples

1) Inducts 2 MilieI Collection of R-modules the product of the {Hi} iEI is the R-modulo $\pi H_i = 2(m_i)_{i \in I} \quad m_i \in H_i f$ With the following universal poperty: given an R-module N and R-module homomorphisms N — Mi there exists a unique R-module homomomorphism f such that $N \xrightarrow{f} \to \pi M_{i}$ commutes for all je I For Ma where π_j : $\pi_i \longrightarrow M_j$ is the projection onto the jth factor Stogan Happing into a product (=> mapping into each factor

(a) Prost sum
$$\{H_{i}\}_{i \in I}$$
 collection \mathcal{Q} R-modules
-the direct sum of the $\{H_{i}\}_{i \in I}$ is the R-module
 $\mathfrak{P}_{i \in I}$ $H_{i} = \{(m_{i})_{i \in I} : m_{i} \in H_{i}, m_{i} = 0 \text{ for all but fittely many } \}$
With the following universal populy:
given an R-module N and R-module homemorphisms $M_{i} = \frac{f_{i}}{N}$
there exists a unique R-module homemorphism such that
 $N \leftarrow f = \mathfrak{P}_{i \in I}$ M_{i}
 $\mathfrak{f}_{i \in I}$ M_{i}
 $\mathfrak{f}_{i \in I}$
where $L_{j}: M_{j} \longrightarrow \mathfrak{P}_{i \in I}$ M_{i} is the inclusion of the jth factor
Stogan Happing out \mathfrak{Q} a direct sum \Leftrightarrow mapping out \mathfrak{Q} each factor
 $\mathfrak{M}_{i \in I}$ $\mathfrak{M}_{i} = \mathfrak{f}_{i \in I}$ $\mathfrak{f}_{i \in I}$ $\mathfrak{f}_{$

Free <u>R-madule</u> the free R-module on $2b_{i}$ fier is the R-module \mp with the following universal property: given an R-module H and a function $f: I \longrightarrow M$, there exists a unique R-module homomorphism $\hat{f}: F \longrightarrow M$ such that



Commutes

$$\mathcal{A}_{j} \quad \hat{f}(b_{i}) = f(i).$$

We say \mp is the free module on I (or $2b_i^2 j_{i \in I}$) and $2b_i^2 j_i^2$ is called a basis.

Stogan to define a map out of a free module is to define any function (of sets) on the basis elements

Real slogen The bosis elements form an actual basis (there are no relations between them, and they generate F)

We saw in the Winter that F≚⊕R = R OLR iEI

Product (category Theory) & category 2 Milier a collection of dojects in G indexed by a set I the product of $2M_i$ lies is the object πM_i together with anows $\pi M_i \xrightarrow{\pi_0} M_j$ with the following universal poperty: given anows $N \xrightarrow{f_i} M_i$ there exits a unique anow f N---> THi Such that Commutes. f_{∂} H_{a} Examples · the product in R-mod is the product of modules. · In Top, the product of topological spaces is the set product with the Tychonoff toplogy (with the bosis of open sets $TT U_i \subseteq TT \times_i$ where is is in the set Un = Xi for all but fritaly money i, all Un open • If I is a gartially ordered set, when we view I as a category we have juduits = greatest bower bounds · In Fredds, there is no Q × Z/p

Coproduct (category theory) & category 2 Milier a collection of dojects in G indexed by a set I the corpoduct of 2Mi Jier is the object II Mi together with anows $M_{j} \xrightarrow{L_{j}} \coprod M_{i}$ with the following universal poperty: given anows $M_i \xrightarrow{f_i} N$ there exists a unique anow f reach that $N \leftarrow \frac{f}{f} - - \prod_{i \in I} H_i$ $f_i \qquad f_i \qquad f_i$ Commutes.

Examples

- In R-mod the copyoduct is the derect sum
- · In Set, the consoluct is the disjoint incom of sets
- · In Gap, the copyeduct is the free groduct

• If I is a gartially ordered set and we orew it as a category, coproducts = least upper bound The dyect (category theory) (concrete category) det \mathcal{C} be a category ruth a faithful functor $F: \mathcal{C} \longrightarrow Set$ det \times be a set

the free edgect on X is an edgect A in Gtogether with an injective function of sets $X \xrightarrow{i} F(A)$ that satisfies the following universal poperty: for any object B in \mathcal{G} and any function $\overline{2}$ sets $x \xrightarrow{f} F(B)$ there exists a unique anow $A \xrightarrow{g} B$ such that $X \xrightarrow{\iota} F(A)$ f $\int F(g)$ F(B)9

Example In R-mod, free modules

When we talk about adjoint functors:

a forgetful functor to set has a left adjoint that is the

Universal peopeting (forwal definition)
G locally small category
A universal peopeting for an depart C in G consists of:
• a superentiable functor F: G
$$\rightarrow$$
 Set
• a universal element $\times \in F(C)$ such that
F is naturally nomorphic to Homge (C, -)
the natural iso conservating to \times by the Homeda hyperten
Escande the universal peoperty of the Ardust $X_4 \times X_2$
is encoded in
• the representable functor $\operatorname{Hom}_{G \times G} (\Delta(-), (X_1, X_2))$
• represented by $X_1 \times X_2$
• tra $(\pi_1, \pi_2) \in \operatorname{Hom}_{G \times G} (\Delta(X_1 \times X_2), (X_1, X_2))$
 $\Delta: \mathcal{E} \longrightarrow \mathcal{E} \times \mathcal{E}$
where $\times (X, X)$
 $f \downarrow \qquad I(f, f)$

•

Unwropping this, it says:
Hom
$$g_{2}(-, x_{1} \times x_{2}) \cong Hom_{g_{2}} g_{2}(\Delta(-), (x_{1}, x_{2}))$$

by the ratival isomorphism γ given by
Hom $g_{2}(\gamma, x_{1} \times x_{2}) \xrightarrow{\varphi_{2}} Hom_{g_{2}} g_{2}(\Delta(\gamma), (x_{1}, x_{2}))$
 $f \longrightarrow \Delta(f)^{*}(\pi, \pi_{2})$
 $id_{\gamma} (\pi_{1}, \pi_{2}) \circ \Delta(f)$
 $\varphi_{\gamma}(f) = (g_{1}, g_{2})$ means
 $x_{1} \xrightarrow{f_{1}} y_{2} \xrightarrow{\chi_{2}} x_{2}$ Commutes
 φ_{γ} is a hypoten
 \Rightarrow for every g_{1}, g_{2} there exists a unique f
rsuch that the diagram commutes.
(that's the universed property of the poduct!)

Hom_R
$$(H_{J}-)$$

and are additive functors $R-mod \longrightarrow R-mod$
Hom_R $(-, N)$

Home
$$(H, \pi, N) \cong \Phi$$
 Home (H_i, N_i)
Home $(H_i, N) \cong \Phi$ Home (H_i, N_i)

• T is left exact if

$$0 \rightarrow A \rightarrow B \rightarrow C$$
 exact $\Rightarrow 0 \rightarrow T(A) \rightarrow T(B) \rightarrow T(C)$ exact
• T is night exact if
 $A \rightarrow B \rightarrow C \rightarrow 0$ exact $\Rightarrow T(A) \rightarrow T(B) \rightarrow T(C) \rightarrow 0$ exact
• T is exact if
 $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0$ hes $\Rightarrow 0 \rightarrow T(A) \rightarrow T(B) \rightarrow T(C) \rightarrow 0$

 $T: R - mod \longrightarrow S - mod \quad (entravariant additive functor)$ T is <u>left exact</u> if $A \rightarrow B \rightarrow C \rightarrow 0 exact \Rightarrow 0 \rightarrow T(C) \rightarrow T(B) \rightarrow T(A) exact$ T is <u>sught exact</u> if $0 \rightarrow A \rightarrow B \rightarrow C \Rightarrow T(C) \rightarrow T(B) \rightarrow T(A) \rightarrow 0 exact$ T is <u>exact</u> if $0 \rightarrow T(C) \rightarrow T(B) \rightarrow T(A) \rightarrow 0 exact$

Remark : Equivalently, ne can start with a ses in all the conditions

- deft exact = turns kernels into kernels (cokernels)
- Right exact = turns cohernels into cohernals (hernels)
- Exact = deft exact + sught exact (contravariant)

<u>theorem</u> (Hom is left exact) • For every short exact sequence $0 \rightarrow A \stackrel{f}{\longrightarrow} B \stackrel{g}{\rightarrow} C \rightarrow 0$ and every R-module M, $0 \rightarrow Hom_{R}(M, A) \stackrel{f_{*}}{\longrightarrow} Hom_{R}(M, B) \stackrel{g_{*}}{\longrightarrow} Hom_{R}(M, C)$ is exact

Theref Concurant version:
(1)
$$f_{\star}$$
 is injective
det $h \in Hom_{R}(M, A)$.
 $f_{\star}(h) = 0 \iff f \circ h = 0 \implies h = 0$
 $f mono$

- (2) Im $f_* \subseteq \ker g_*$ For all he Hom_R(H,A), $g_*f_*(h) = g_{\circ}f_{\circ}h = 0$ (3) $\ker g_* \subseteq \operatorname{Im} f_*$ det he Hom_R(H,B), $g_*(h) = 0 \iff g_{\circ}h = 0$ For all meth, $g(h(m)) = 0 \Rightarrow h(m) \in \ker g = \operatorname{Im} f$ (hoose a $\in A$ such that f(a) = h(m), and set $\kappa(m) := a$
- $\int ungerture \implies for each m \exists unique a \implies k is nell-defined$
- Exercise: $k \in Hom_{\mathcal{R}}(M, A)$ defoza
- $f_{*}(k)$ satisfies $f_{*}(k)(m) = f(k(m)) = f(a) \stackrel{l}{=} h(m)$ $\Rightarrow f_{*}(k) = h$.

Example
$$0 \rightarrow Z \xrightarrow{i} Q \xrightarrow{\pi} Q/Z \rightarrow 0$$

 $\downarrow \operatorname{Hom}_{Z}(Z_{2}, -)$
 $0 \rightarrow \operatorname{Hom}_{Z}(Z_{2}, Z) \xrightarrow{i_{w}} \operatorname{Hom}_{Z}(Z_{2}, Q) \xrightarrow{\pi_{w}} \operatorname{Hom}_{Z}(Z_{2}, Q)$
 $Claim \quad \pi_{x} \quad \operatorname{not} \quad \operatorname{suggettre}$
 $\cdot \operatorname{Hom}_{Z}(Z_{2}, Q) = 0 \quad (1 \mapsto a, \quad 2a = 0 \Rightarrow a = 0)$
 $\cdot \operatorname{Hom}_{Z}(Z_{2}, Q) = 0 \quad (1 \mapsto a, \quad 2a = 0 \Rightarrow a = 0)$
 $\cdot \operatorname{Hom}_{Z}(Z_{2}, Q/Z) \neq 0$
 $1 \mapsto \frac{1}{2} + Z \quad \text{is de !}$
so actually, that exact sequence is
 $0 \rightarrow 0 \rightarrow 0 \rightarrow Hom_{Z}(2/2, Q/Z)$
 $\xrightarrow{\omega} Z/Q \qquad \checkmark$
Also, when we apply $\operatorname{Hom}_{Z}(-, Z)$, we get
 $0 \rightarrow \operatorname{Hom}_{Z}(Q/Z, Z) \xrightarrow{g^{*}} \operatorname{Hom}_{Z}(Q, Z) \xrightarrow{f^{*}} \operatorname{Hom}_{Z}(Z, Z)$
 $\Rightarrow f^{*} \text{ net suggettre !}$