

# CHAMPS 2021

Minimal DG Algebra Resolutions of Certain Stanley-Reisner Rings

## GOALS

- ① Convince you of the value of  $\hat{\Delta}$
- ② Showcase a minimal free resolution
- ③ Showcase the DG algebra structure

Goal 1 :  $\Delta$

Definition: A simplicial complex  $\Delta$  on a vertex set  $V$  is a collection of subsets of  $V$  which is closed under taking subsets. The elements of  $\Delta$  are called the faces of the complex and the maximal faces are the facets.

Example:  $V = \{a, b, c\}$

$$\Delta = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, c\}\} = \{\emptyset, a, b, c, ac\} = \langle ac, b \rangle = \left( \begin{array}{c} c \\ a \cdot \text{---} \cdot c \\ \cdot b \end{array} \right)$$

Definition: A simplicial complex is pure if all its facets are the same size. If  $\Delta$  is a simplicial complex on  $V = \{a_1, \dots, a_n\}$ , then let  $\hat{V} = V \cup \{\alpha, \dots, \alpha_n\}$  be a vertex set and for each face  $F \in \Delta$  define

$$\hat{F} = F \cup \{\alpha_i \in \hat{V} \mid a_i \notin F\}.$$

Let  $\hat{\Delta}$  be the simplicial complex on  $\hat{V}$  generated by all such  $\hat{F}$  and call it the purification of  $\Delta$ .

Example:  $\hat{V} = V \cup \{\alpha, \beta, \gamma\}$ ,  $\Delta = \{\emptyset, a, b, c, ac\}$

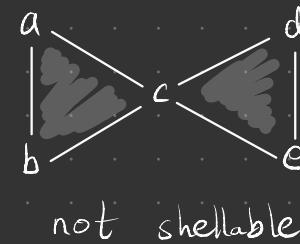
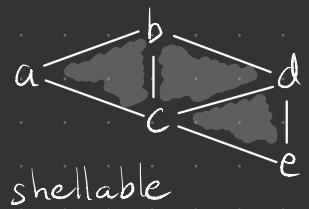
<u><math>F</math></u>	<u><math>F^c \subset V</math></u>	<u>Convert to Greeks</u>	<u>combine</u>
$\emptyset$	$abc$	$\alpha\beta\gamma$	$\alpha\beta\gamma$
$a$	$bc$	$\beta\gamma$	$a\beta\gamma$
$ac$	$b$	$\beta$	$a\beta c$

$$\therefore \hat{\Delta} = \langle \alpha\beta\gamma, a\beta\gamma, \alpha b\gamma, \alpha\beta c, a\beta c \rangle = \left( \begin{array}{c} \text{Diagram of a tetrahedron with vertices } a, b, c \text{ and faces } \alpha, \beta, \gamma. \end{array} \right)$$

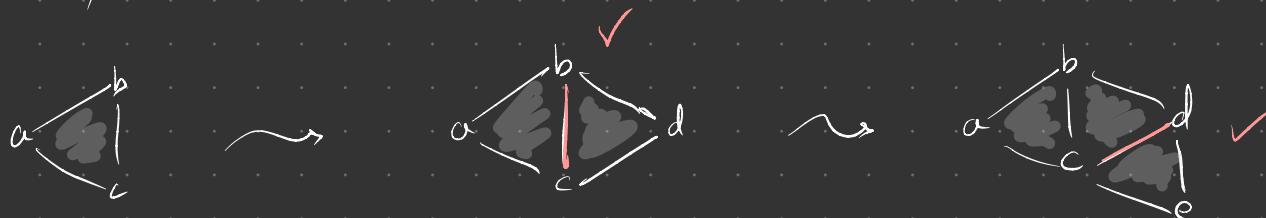
Definition: a pure simplicial complex is shellable if one can list its facets

$F_1, F_2, \dots, F_p$  such that  $\dim(\langle F_i \rangle \cap \langle F_1, \dots, F_{i-1} \rangle) = \dim(F_i) - 1$  for every  $i=2, \dots, p$ .

Example:



shelling: abc, bed, cde



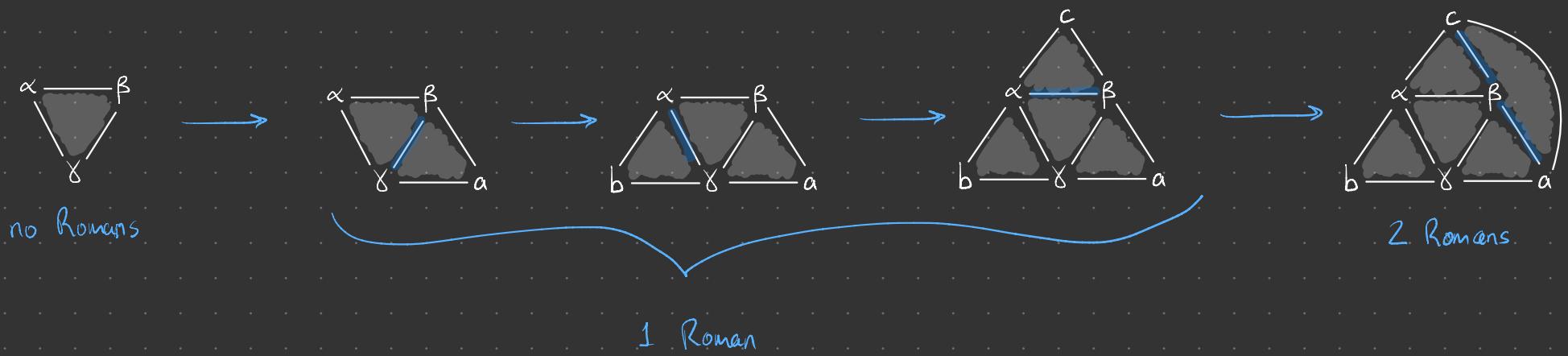
not a shelling: abc, cde, bcd



Goal 1 :  $\hat{\Delta}$

Theorem (Morra)  $\hat{\Delta}$  is shellable. Specifically, if the facets of  $\hat{\Delta}$  are listed in order of increasing number of Romans, then that list is a shelling.

Example:  $\alpha\beta\gamma, \alpha\beta\delta, \alpha\delta\gamma, \alpha\beta\gamma, \alpha\beta\gamma$  is a shelling



Remark: shellable  $\Rightarrow$  C-M [Bruns & Herzog, Theorem 5.1.3]

Goal 1 :  $\hat{\Delta}$

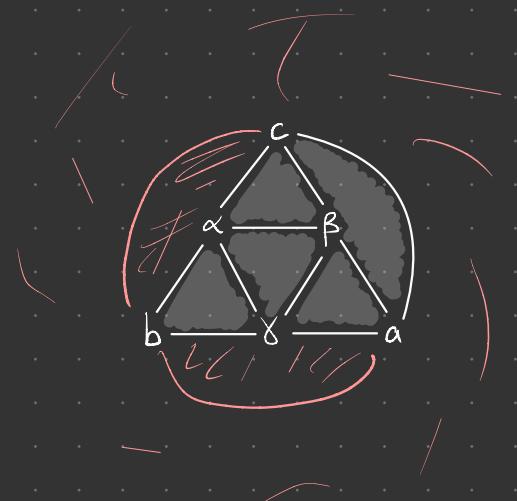
Definition: Let  $\Delta$  and  $\hat{\Delta}$  be as above, identify a variable with each vertex in  $\hat{V}$ , and define the polynomial ring  $R = k[a_1, \dots, a_n, \alpha_1, \dots, \alpha_n]$ . The Stanley-Reisner ideal of  $\hat{\Delta}$  is generated by the non-faces of  $\hat{\Delta}$ , i.e.,

$$J_{\hat{\Delta}} = \langle m_F \mid F \subset \hat{V}, F \notin \hat{\Delta} \rangle \subseteq R$$

where  $m_F = \prod_{v \in F} v$ . The Stanley-Reisner ring of  $\hat{\Delta}$  is  $S = R/J_{\hat{\Delta}}$ .

Example: If  $\hat{\Delta} = \langle \alpha\beta\gamma, \alpha\beta\delta, \alpha b\delta, \alpha\beta c, \alpha b c \rangle$ , then

$$\begin{aligned} J_{\hat{\Delta}} &= \langle \alpha\alpha, b\beta, c\gamma, ab, bc, ab\delta, \alpha bc, abc, \dots \rangle \\ &= \langle \alpha\alpha, b\beta, c\gamma, ab, bc \rangle \\ &= \langle a, b, c \rangle \cap \langle \alpha, b, c \rangle \cap \langle \alpha, \beta, c \rangle \cap \langle a, b, \delta \rangle \cap \langle \alpha, b, \delta \rangle \end{aligned}$$



Remark:  $J_{\hat{\Delta}}$  is unmixed.

Theorem (Morra): The minimal generators of the canonical module  $\omega_S$  of  $S$  are in bijection with the facets of  $\Delta$ .

Example:  $\Delta = \langle b, ac \rangle$        $\hat{\Delta} = \langle \alpha\beta\gamma, \alpha\beta\delta, \alpha b\delta, \alpha\beta c, \alpha\beta c \rangle$

$$\begin{array}{ccccc} \{b\} & \xrightarrow{\quad} & \{a, c\} & \xrightarrow{\quad} & \{\alpha, \gamma\} \\ & \text{complement} & & \text{change to Greek} & \\ \{a, c\} & \xrightarrow{\quad} & \{b\} & \xrightarrow{\quad} & \{\beta\} \end{array} \quad \left. \begin{array}{c} \nearrow \\ \searrow \end{array} \right\} \omega_S = \{\alpha\gamma, \beta\} \subseteq S$$

Fact: If we localize  $S$ , e.g.,  $S_*$  where  $* \subseteq S$  is generated by the variables, then the Cohen-Macaulay type of  $S_*$  is equal to the number of facets in  $\Delta$ .

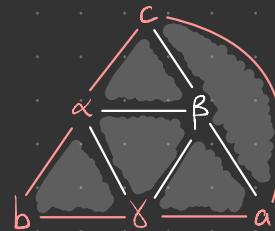
Corollary: In the special case when  $J_{\hat{\Delta}}$  is the edge ideal of a  $K_1$ -corona  $\Sigma G$  of a simple graph  $G$ , the Cohen-Macaulay type of  $S_*$  is equal to the number of maximal cliques in  $G^c$ .

Fact: Excluding the case when  $\Delta$  is a simplex,  $\hat{\Delta}$  is topologically equivalent to a ball, so its boundary is a sphere. [D'ali, Fløystad, Nematbakhsh (2019)]

Example: Let  $\Sigma$  denote the boundary of  $\hat{\Delta}$ .

$$\hat{\Delta} = \langle \alpha\beta\gamma, \alpha\beta\delta, \alpha\delta\gamma, \alpha\beta\gamma, \alpha\beta\gamma \rangle$$

$$\Sigma = \langle ac, ab, b\gamma, \alpha b, \alpha c \rangle$$



$$\hat{\Delta} \setminus \Sigma = \{ \beta, \alpha\beta, \alpha\delta, \beta\delta, \alpha\beta, \beta\gamma, \alpha\beta\gamma, \alpha\beta\delta, \alpha\delta\gamma, \alpha\beta\gamma, \alpha\beta\gamma \}$$

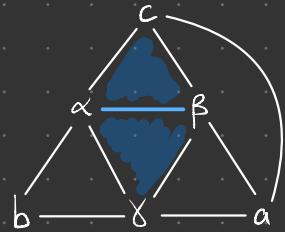
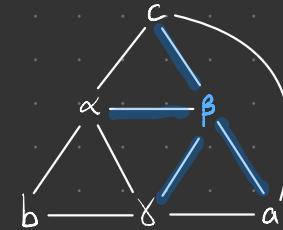
[      ] [      ]  
 $B_3$              $B_2$              $B_1$

## Goal 2: Resolution

Example: Let  $\mathcal{L}$  denote the following resolution:

$$\begin{array}{ccccccc}
 S & \xleftarrow{\text{0}} & S^{(B_1)} & \xleftarrow{\text{1}} & S^{(B_2)} & \xleftarrow{\text{2}} & S^{(B_3)} & \xleftarrow{\text{3}} & O \\
 & & & & & & & & \\
 & & [\alpha\beta\gamma] & [\alpha\beta\gamma] & [\alpha\beta\gamma] & [\alpha\beta] & [\beta] \\
 & & [\alpha\beta c] & [\alpha\beta c] & & [\alpha\beta] & [\beta c] \\
 & & & & & & 
 \end{array}$$

$$\delta[\beta] = -\alpha[\alpha\beta] - \alpha[\alpha\beta] + \gamma[\beta\gamma] + c[\beta c]$$



$$\delta[\alpha\beta] = \gamma[\alpha\beta\gamma] + c[\alpha\beta c]$$

$$\delta[\alpha\beta\gamma] = abc$$

$$\delta[\alpha\beta\gamma] = -\alpha bc$$

$$\delta[\alpha\beta c] = \alpha b\gamma$$

Goal 2: Resolution

$\mathcal{L}$  minimally resolves  $S/\mathcal{I}$ , where  $\mathcal{I} = (\mathcal{J}_{\Delta})^A$ .

$$\begin{array}{ccccccc}
 S & \xleftarrow{\text{1}} & S^{(B_1)} & \xleftarrow{\text{2}} & S^{(B_2)} & \xleftarrow{\text{3}} & S^{(B_3)} \xleftarrow{} O \\
 [abc] & [abc] & [abc] & & [\alpha\beta\gamma] & [\alpha\gamma] & [\beta\gamma] \\
 & [\alpha\beta c] & [\alpha\beta c] & & [\alpha\beta] & [\beta c] & [\beta]
 \end{array}$$

$$\mathcal{J}_{\Delta} = \langle a, b, c \rangle \cap \langle \alpha, b, c \rangle \cap \langle a, \beta, c \rangle \cap \langle a, b, \gamma \rangle \cap \langle \alpha, b, \gamma \rangle$$

$$\Rightarrow \mathcal{I} = (\mathcal{J}_{\Delta})^A = \langle abc, \alpha bc, \alpha\beta c, ab\gamma, \alpha b\gamma \rangle$$

Fact: By [DFN, Theorem 3.11], this is always a minimal resolution.

Definition: A commutative differential graded S-algebra is an S-complex

$$A = \left( 0 \leftarrow A_0 \leftarrow A_1 \leftarrow A_2 \leftarrow \dots \right) \right)$$

equipped with a binary operator  $\cdot$  such that  $|a \cdot b| = |a| + |b|$  and it

- (i) is S-bilinear,
- (ii) is unital, i.e.,  $\exists 1 \in A_0$  st.  $1 \cdot a = a = a \cdot 1$ ,  $\forall a \in A$ ,
- (iii) is associative,
- (iv) is graded commutative, i.e.,
  - o  $\forall a, b \in A \setminus \{0\}$  one has  $b \cdot a = (-1)^{|a||b|} a \cdot b$ , and
  - o  $|a| \text{ odd } \Rightarrow a \cdot a = 0$
- (v) satisfies the Leibniz rule, i.e.,  $\forall a, b \in A \setminus \{0\}$  one has

$$\delta(a \cdot b) = \delta(a) \cdot b + (-1)^{|a|} a \cdot \delta(b).$$

Example: Koszul & Taylor.

Example:

$$\begin{array}{ccccccc}
 S & \xleftarrow{\quad 0 \quad} & S^{(\mathcal{B}_1)} & \xleftarrow{\quad 1 \quad} & S^{(\mathcal{B}_2)} & \xleftarrow{\quad 2 \quad} & S^{(\mathcal{B}_3)} & \xleftarrow{\quad 3 \quad} & 0 \\
 & & [\alpha\beta\delta] & [\alpha\beta\gamma] & [\alpha\delta] & [\alpha\gamma] & [\beta\delta] & [\beta\gamma] & \\
 & & [\alpha\beta c] & [\alpha\gamma c] & & [\alpha\beta] & [\beta c] & 
 \end{array}$$

$$[\alpha\beta\delta][\alpha\beta\gamma] = ?$$

$$|[\alpha\beta\gamma][\alpha\beta\delta]| = |[\alpha\beta\delta]| + |[\alpha\beta\gamma]| = 2$$

$$\therefore [\alpha\beta\delta][\alpha\beta\gamma] = s_1[\alpha\beta] + s_2[\alpha\gamma] + s_3[\beta\delta] + s_4[\alpha\delta] + s_5[\beta\gamma]$$

$$\therefore \partial([\alpha\beta\delta][\alpha\beta\gamma]) = s_1 \cdot \partial([\alpha\beta]) + s_2 \cdot \partial([\alpha\gamma]) + s_3 \cdot \partial([\beta\delta]) + s_4 \cdot \partial([\alpha\delta]) + s_5 \cdot \partial([\beta\gamma]) \quad (1)$$

Leibniz:

$$\begin{aligned}
 \partial([\alpha\beta\delta][\alpha\beta\gamma]) &= \partial([\alpha\beta\delta]) \cdot [\alpha\beta\gamma] + (-1)^1 [\alpha\beta\delta] \cdot \partial([\alpha\beta\gamma]) \\
 &= abc[\alpha\beta\delta] - abc[\alpha\beta\gamma]
 \end{aligned} \quad (2)$$

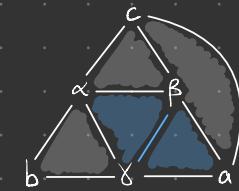
Much bookkeeping later:

$$[\alpha\beta\delta][\alpha\beta\gamma] = bc[\beta\delta].$$

### Goal 3: DG algebra

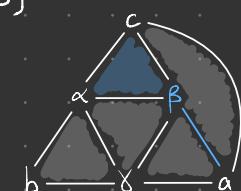
Example:  $[\alpha\beta\gamma][\alpha\beta\gamma] = bc[\beta\gamma]$      $\alpha\beta\gamma \cap \alpha\beta\gamma = \beta\gamma \in \hat{\Delta} \setminus \Sigma$  and  $|[\beta\gamma]| = 2$  ✓

graded commutativity  $\Rightarrow [\alpha\beta\gamma][\alpha\beta\gamma] = -(-1)^{1+1} bc[\beta\gamma] = -bc[\beta\gamma]$



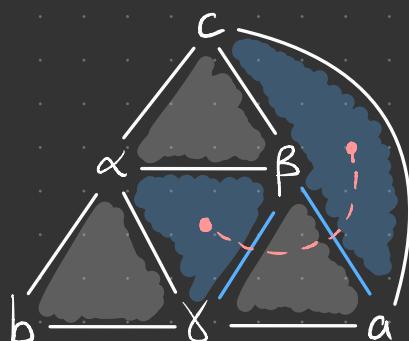
Example:  $\alpha\beta c \cap \alpha\beta = \beta$  and  $|[\alpha\beta c]| + |[\alpha\beta]| = 1 + 2 = |[\beta]|$   $\therefore [\alpha\beta c][\alpha\beta] = b\gamma[\beta]$

graded commutativity  $\Rightarrow [\alpha\beta][\alpha\beta c] = (-1)^{1+2} b\gamma[\beta] = b\gamma[\beta]$



Example:  $[\alpha\beta\gamma][\alpha\beta c] = -b\gamma[\beta\gamma] + ab[\alpha\beta]$

$\alpha\beta\gamma \cap \alpha\beta c = \gamma \in \hat{\Delta} \setminus \Sigma$  ✓     $|[\beta]| = 3 \neq |[\alpha\beta\gamma]| + |[\alpha\beta c]|$  ✗



$\alpha\beta\gamma, \alpha\beta\gamma, \alpha\beta$

### Goal 3: DG algebra

#### Discussion.

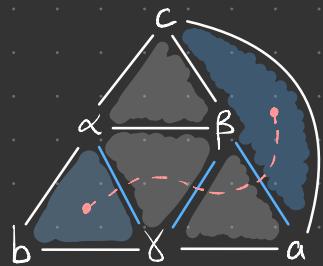
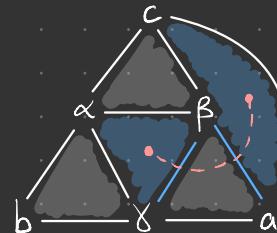
① When  $F \cap F'$  is of appropriate codimension and is an element of  $\hat{\Delta} \setminus \Sigma$ ,

the product  $[F][F']$  is simple.  $\pm (F^c \cap F'^c) [F' \cap F]$

② When  $F \cap F'$  has the wrong codimension, the product  $[F][F']$  is a linear combination of simple products that arise from a "path from  $F$  to  $F'$ ."  
Such products we call complex.

$$[\alpha\beta\gamma][a\beta c] = (-1)^3 \left( \frac{\gamma}{c} [\alpha\beta\gamma][a\beta\gamma] + \frac{a}{\alpha} [\alpha\beta\gamma][a\beta c] \right)$$

$$= -b\gamma[\beta\gamma] + ab[\alpha\beta]$$



$$[\alpha\beta\gamma][a\beta c] = (-1)^4 \left( \frac{\alpha\gamma}{ac} [\alpha\beta\gamma][\alpha\beta\gamma] + \frac{\beta\gamma}{bc} [\alpha\beta\gamma][a\beta\gamma] + \frac{\alpha\beta}{ab} [\alpha\beta\gamma][a\beta c] \right)$$

$$= -\alpha\gamma[\alpha\gamma] + \beta\gamma[\beta\gamma] - \alpha\beta[a\beta]$$

Corollary:  $\mathcal{I}$  is a Golod ideal.

### Goal 3: DG algebra

Definition: For each  $i=1, \dots, n$  define  $\tau_i: \hat{V} \rightarrow \hat{V}$  by

$$\tau_i(F) = (F \setminus \{a_i\}) \cup \{\alpha_i\}$$

if  $a_i \in F$ . Similarly, we define  $t_i: \hat{V} \rightarrow \hat{V}$  by

$$t_i(F) = (F \setminus \{\alpha_i\}) \cup \{a_i\}$$

if  $\alpha_i \in F$ .

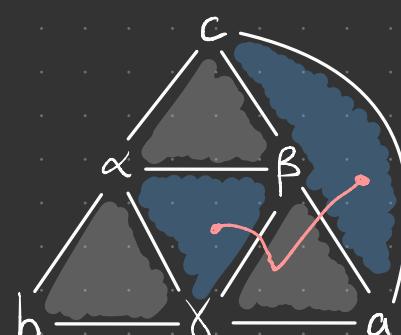
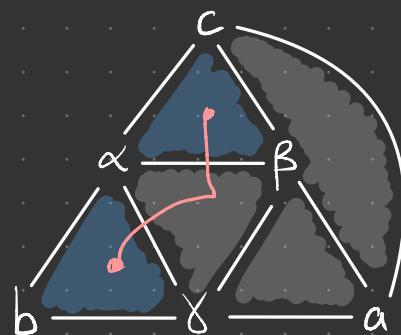
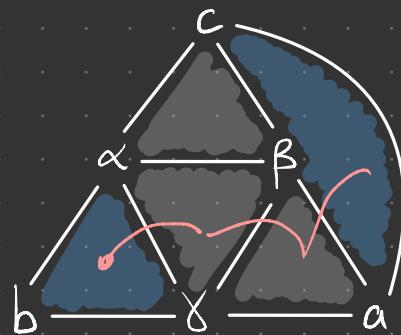
Example:  $\alpha b \gamma \xrightarrow{\tau_2} \alpha \beta \gamma \xrightarrow{t_1} \alpha \beta \gamma \xrightarrow{t_3} \alpha \beta c$

$$P(\alpha b \gamma, \alpha \beta c) = \{\alpha b \gamma, \alpha \beta \gamma, \alpha \beta \gamma, \alpha \beta c\}$$

"path from  $\alpha b \gamma$  to  $\alpha \beta c$ "

Remark: When forming paths between facets, we apply these maps as-needed in this order:

$$\tau_n, \tau_{n-1}, \dots, \tau_1, t_1, t_2, \dots, t_n.$$



### Goal 3: DG algebra

Discussion. Many products are zero under this structure,

including some when the intersection is of appropriate codimension, e.g.,

$$[\alpha\beta c][a\beta] = b\gamma[\beta], \quad \text{but} \quad [\alpha\beta c][\beta\gamma] = 0,$$

even though we have

$$|[\alpha\beta c \cap a\beta]| = |[\beta]| = 3 \quad |[\alpha\beta c \cap \beta\gamma]| = |[\beta]| = 3.$$

It turns out that for each face  $F \in \hat{\Delta} \setminus \Sigma$  we can define a set  $\varepsilon(F) \subset \Phi(\hat{V})$  such that under a few reasonable conditions we have

$$[F][F'] \neq 0 \quad \text{iff} \quad F' \in \hat{\Delta} \setminus \Sigma \text{ contains an element of } \varepsilon(F).$$

Example:  $\varepsilon(\alpha\beta c) = \{\alpha\beta\gamma, a\beta, b\} \subset \Phi(\hat{V})$

$$\therefore [\alpha\beta c][a\beta] \neq 0 \quad \text{and} \quad [\alpha\beta c][\beta\gamma] = 0$$

Example:  $\varepsilon(a\beta) = ?$ : a bit tedious

Thank You.